Optimization and tensor decompositions

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Why do we need data models?

An array $A = [a(i_1, \ldots, i_d)]$ of size $n \times \ldots \times n$ cannot be given by the list of all entries even for modest values of d.

If d = 83 and n = 10, then the number of entries 10^{83} equals the number of atoms in the Universe!



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We need *data representation models* with small sets of parameters and *algorithms dealing only with these small sets*.

A data-model paradigm of computations

- "Big objects" are represented through "small vectors" A = A(p), B = B(q), C = C(s).
- Computing a "big object" C = A * B is done by a *fast algorithm* that finds s from p and q. "Big objects" A and B never arise explicitly.
- Data models for a wide class of applications can be based on low-rank matrices.

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Low-rank matrix model

If A is a matrix of rank r, then

$$A = \sum_{\alpha=1}^{r} u_{\alpha} v_{\alpha}^{\top} = UV^{\top}$$

$$U = [u_1, \ldots, u_r], \quad V = [v_1, \ldots, v_r]$$

(matrices of size $n \times r$)

 $2rn \ll n^2$

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Advantages of low rank

Data compression

► for matrices of size $n \times n$: $r \ll n \Rightarrow 2nr \ll n^2$

► for *d*-tensors of size $n \times ... \times n$: $r \ll n^{d-1} \Rightarrow dnr \ll n^d$

- Extraction of most meaningful information and suppression of noise
- Fast computations in low-rank formats

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Singular Value Decomposition (SVD)

It is a skeleton decomposition in which vectors u_1, \ldots, u_r unterpretent v_1, \ldots, v_r are orthogonal..

Afer normalization

$$u_{\alpha}v_{\alpha}^{\top} = \sigma_{\alpha}u_{\alpha}v_{\alpha}^{\top}, \quad u_{\alpha} := u_{\alpha}/||u_{\alpha}||, \quad v_{\alpha} := v_{\alpha}/||v_{\alpha}||,$$

we come up with a *Singular Value Decomposition*:

$$A = \sum_{\alpha=1}^{r} \sigma_{\alpha} u_{\alpha} v_{\alpha}^{\top} = U \Sigma V^{\top}, \quad \Sigma = \begin{bmatrix} \sigma_{1} & & \\ & \ddots & \\ & & \sigma_{r} \end{bmatrix},$$

$$\sigma_1 \geqslant \sigma_2 \geqslant \ldots \geqslant \sigma_r > \sigma_{r+1} = \ldots = \sigma_n = 0.$$

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Low-rank approximations by a truncated SVD

$$\min_{\operatorname{rank} B \leqslant k} ||A - B|| = ||A - A_k|| = \sqrt{\sum_{\alpha=k+1}^n \sigma_{\alpha}^2}$$

$$A_k = \sum_{\alpha=1}^k \sigma_\alpha u_\alpha v_\alpha^\top$$

$$||A|| = \sqrt{\sum_{i,j} |a_{ij}|^2}$$
 (Frobenius norm)

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Tensor Train Model

$$x_{i_1,\ldots,i_d} = \sum_{\alpha_1=1}^{r_1} \cdots \sum_{\alpha_{d-1}=1}^{r_{d-1}} g_{i_1,\alpha_1}^{(1)} g_{\alpha_1,i_2,\alpha_2}^{(2)} \cdots g_{\alpha_{d-2},i_{d-1},\alpha_{d-1}}^{(d-1)} g_{\alpha_{d-1},i_d}^{(d)}$$

Symbolic notation: $x = g^{(1)} \dots g^{(d)}$

Total number of TT-model parameters depends on *d* linearly: $\sum_{k=1}^{d} r_{k-1}r_kn_k \leq dnr^2$, $n = \max_k n_k$, $r = \max_k r_k$.

The train car sizes r_k are called TT-ranks.

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Options to represent the train cars

▶ by matrices of size of size $r_{k-1} \times r_k$: $(G_{i_k}^{(k)})_{\alpha_{k-1},\alpha_k} := g_{\alpha_{k-1},i_k,\alpha_k}^{(k)}$

$$x_{i_1,\ldots,x_d} = G_{i_1}^{(1)}G_{i_2}^{(2)}\ldots G_{i_{d-1}}^{(d-1)}G_{i_d}^{(d)},$$

▶ by column vectors of size n_k : $(g^{(k)}(\alpha_{k-1}, \alpha_k))_{i_k} := g^{(k)}_{\alpha_{k-1}, i_k, \alpha_k}$

$$x = \sum_{\alpha_1=1}^{r_1} \dots \sum_{\alpha_{d-1}=1}^{r_{d-1}} g^{(1)}(\alpha_0, \alpha_1) \otimes g^{(2)}(\alpha_1, \alpha_2) \otimes \dots \otimes g^{(d)}(\alpha_{d-1}, \alpha_d).$$

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$r_k \ge \operatorname{rank}(A_k), \quad (A_k)_{(i_1,\ldots,i_k),(i_{k+1},\ldots,i_d)} := x_{i_1,\ldots,i_d}.$

Theorem.

Any tensor admits a TT-representation with $r_k = \operatorname{rank}(A_k)$.

Find an approximation

$$\tilde{x} = \tilde{g}^{(1)} \dots \tilde{g}^{(d)} \approx x = g^{(1)} \dots g^{(d)}$$

with smaller TT-ranks and a guaranteed accuracy.

Reducing the rank of kth car

$$\begin{aligned} x_{i_1,\dots,i_d} &= G_{i_1}^{(1)} \dots G_{i_d}^{(d)} \\ u_{i_1,\dots,i_k} &= G_{i_1}^{(1)} \dots G_{i_k}^{(k)} \ \text{(rows)}, \quad v_{i_{k+1},\dots,i_d} = G_{i_{k+1}}^{(k+1)} \dots G_{i_d}^{(d)} \ \text{(columns)} \\ A_k &= U_k V_k \end{aligned}$$

• Orthogonalize columns in U_k and rows in V_k :

$$U_k = P_k S_k, \quad dV_k = T_k Q_k \quad \Rightarrow \quad A_k = U_k V_k = P_k B_k Q_k, \quad B_k = S_k T_k.$$

- Approximate $B_k \approx \tilde{B}_k$ with $\tilde{r}_k < r_k$ and $||B_k \tilde{B}_k|| = \varepsilon$.
- Define tensor \tilde{x} by $\tilde{A}_k := P_k \tilde{B}_k Q_k$. Then

$$||x-\tilde{x}|| = ||A_k - \tilde{A}_k|| = ||B_k - \tilde{B}_k|| = \varepsilon.$$

Rank reduction can be fast

$$\Pi'_k = \begin{bmatrix} G_1^{(k)} \\ \cdots \\ G_{n_k}^{(k)} \end{bmatrix}, \quad \Pi''_k = \begin{bmatrix} G_1^{(k)} & \cdots & G_{n_k}^{(k)} \end{bmatrix}.$$

- ► Train car $g^{(k)}$ is called *left-orthogonal* if Π'_k has orthonormal columns.
- Train car $g^{(k)}$ is called *right-orthogonal* if $\Pi_k^{\prime\prime}$ has orthonormal rows.

- ▶ If $g^{(1)}, \ldots, g^{(k)}$ are left-orthogonal, then U_k has orthonormal columns.
- If $g^{(k+1)}, \ldots, g^{(d)}$ are right-orthogonal, then V_k has orthonormal rows.

Make the first car left-orthogonal

$$(U_k)_{(i_1,...,i_k),lpha_k} = \sum_{lpha_1,...,lpha_{k-1}} g_{i_1,lpha_1}^{(1)} \dots g_{lpha_{k-1},i_k,lpha_k}^{(k)}$$

 $g_{i_1,lpha_1}^{(1)} = \sum_{eta_1} p_{i_1,eta_1}^{(1)} s_{eta_1,lpha_1}^{(1)}$
 $g_{eta_1,i_2,lpha_2}^{(2)} := \sum_{lpha_1} s_{eta_1,lpha_1}^{(1)} g_{lpha_1,i_2,lpha_2}^{(2)}$

$$(U_k)_{(i_1,\ldots,i_k),\alpha_k} = \sum_{\beta_1,\alpha_2,\ldots,\alpha_{k-1}} p_{i_1,\beta_1}^{(1)} g_{\beta_1,i_2,\alpha_2}^{(2)} \cdots g_{\alpha_{k-1},i_k,\alpha_k}^{(k)}$$

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Fast structured SVD for the associated matrix A_k

$$(U_k)_{(i_1,\ldots,i_k),lpha_k} = \sum_{eta_1,\ldots,eta_{k-1},eta_k} p_{i_1,eta_1}^{(1)} p_{eta_1,i_2,eta_2}^{(2)} \ldots p_{eta_{k-1},i_k,eta_k}^{(k)} s_{eta_k,lpha_k}^{(k)}$$

 $(V_k)_{lpha_k,(i_{k+1},\ldots,i_d)} = \sum_{\gamma_k,\gamma_{k+1},\ldots,\gamma_{d-1}} t_{lpha_k,\gamma_k}^{(k+1)} q_{\gamma_k,i_{k+1},\gamma_{k+1}}^{(k+1)} \ldots q_{\gamma_{d-1},i_d}^{(d)}$

$$x_{i_1,\ldots,i_d} = \sum_{\beta_1,\ldots,\beta_k,\gamma_k,\ldots,\gamma_{d-1}} p_{i_1,\beta_1}^{(1)} \ldots p_{\beta_{k-1},i_k,\beta_k}^{(k)} h_{\beta_k,\gamma_k}^{(k)} q_{\gamma_k,i_{k+1},\gamma_{k+1}}^{(k+1)} \ldots q_{\gamma_{d-1},i_d}^{(d)}$$

$$h^{(k)}_{eta_k,\gamma_k} \;=\; \sum_{lpha_k} s^{(k)}_{eta_k,lpha_k} \, t^{(k+1)}_{lpha_k,\gamma_k}$$

 $g^{(1)} \ldots g^{(d)} = p^{(1)} \ldots p^{(k)} h^{(k)} q^{(k+1)} \ldots q^{(d)}$

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$g_1g_2g_3g_4g_5 \rightarrow g_1g_2g_3g_4g_5 \rightarrow g_1g_2g_3g_4g_5 \rightarrow g_1g_2g_3g_4g_5$

 $g_1g_2g_3g_4g_5 \rightarrow g_1g_2g_3g_4g_5 \rightarrow g_1g_2g_3g_4g_5 \rightarrow g_1g_2g_3g_4g_5$

FAST: total number of operations is of order *dnr*³. SAFE: accuracy estimate is guaranteed.

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Cross Column-Row Approximation



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How to find a "good cross"



Maximal volume principle

THEOREM (Goreinov, Tyrtyshnikov).

Assume that $A_{11} \in \mathbb{C}^{r \times r}$ is a maximal volume block among all $r \times r$ blocks in

$$A = \begin{bmatrix} A_{11} & A_{12} \\ A_{21} & A_{22} \end{bmatrix}, \quad r = \operatorname{rank} (A + F).$$

Then

$$||A - CA_{11}^{-1}R||_C \leq (r+1)^2 ||F||_C,$$

 $C = \begin{bmatrix} A_{11} \\ A_{21} \end{bmatrix}, \quad R = \begin{bmatrix} A_{11} & A_{12} \end{bmatrix}.$

Using more columns and rows

Can we obtain a better rank-r approximation using more columns and rows? And how much better?

Theorem (N.Zamarashkin, A.Osinsky'2017).

Assume that $A_{11} \in \mathbb{C}^{p \times q}$ is of rank not smaller than r and has the maximal r-projective volume

$$\mathcal{V}(A_{11}) := \prod_{i=1}^r \sigma_i(A)$$

among all $p \times q$ submatrices in A. Then

$$||A - C(A_{11})_r^{\dagger} R||_C \leq \sqrt{1 + \frac{r}{p - r + 1}} \sqrt{1 + \frac{r}{q - r + 1}} ||F||_2.$$

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How to find a "good volume block"

Pick up an invertible block $\hat{C} \in \mathbb{R}^{k \times k}$ in $C \in \mathbb{R}^{n \times k}$ and get to a matrix

$$C\hat{C}^{-1} = egin{bmatrix} 1 & & & \ & \ddots & & \ & & 1 \ q_{r+1,1} & \dots & q_{r+1,r} \ & \dots & \dots & \dots \ q_{n1} & \dots & q_{nr} \end{bmatrix}$$

For \hat{C} to be of maximal volume it is necessary that $|q_{ij}| \leq 1$, $r+1 \leq i \leq n$, $1 \leq j \leq r$. If not, then swap two rows and increase the volume!

D.Knuth, Semi-optimal bases for linear dependencies, Linear and Mulilinear Algebra, 1985



In the column submatrix $C \in \mathbb{R}^{n \times k}$ look for k "good rows". They determine a row submatrix $R \in \mathbb{R}^{k \times n}$ where we select k "good columns". The new columns give us a new columns submatrix $C \in \mathbb{R}^{k \times n}$ in which we pick up k "good rows". And so on.

With high probability this algorithm leads to a block of close to maximal volume. It becomes a base for a new heuristic global optimization algorithm. For some applicaions (eg docking in the drug design) it appears 2-3 orders better than genetic algorithms.

Direct docking in the drug design

ACCOMMODATION OF LIGAND INTO PROTEIN



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Direct docking in the drug design

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Tensor Train Learning and Optimization

Tensor Train:

$$\begin{aligned} a_{i_1,\ldots,i_d} &= \sum_{\alpha_1=1}^{r_1} \cdots \sum_{\alpha_{d-1}=1}^{r_{d-1}} g_{i_1\alpha_1}^{(1)} g_{\alpha_1 i_2 \alpha_2}^{(2)} \cdots g_{\alpha_{d-2} i_{d-1} \alpha_{d-1}}^{(d-1)} g_{\alpha_{d-1} i_d}^{(d)} \\ &= G_{i_1}^{(1)} G_{i_2}^{(2)} \cdots G_{i_{d-1}}^{(d-1)} G_{i_d}^{(d)} \end{aligned}$$

Algorithms use structured low-rank representations of certain associated matrices

$$A_k = [a_{i_1 \dots i_k; i_{k+1} \dots i_d}^k]_{(n_1 \dots n_k) \times (n_{k+1} \dots n_d)}$$

$$a_{i_1\ldots i_k;i_{k+1}\ldots i_d}^k = a_{i_1,\ldots,i_d}$$

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Fast summation of astronomically big vector

$$i = \overline{i_1 i_2 \dots i_d}$$
 $d = 83$

$$a(i) = a(i_1, \ldots, i_d) = \sum_{\alpha_1, \ldots, \alpha_{d-1}} g_1(i_1, \alpha_1) g_2(\alpha_1, i_2, \alpha_2) \ldots g_d(\alpha_{d-1}, i_d)$$

$$\sum_{i_1,\ldots,i_d} a(i_1,\ldots,i_d) = \sum_{\alpha_1,\ldots,\alpha_{d-1}} \hat{g}_1(\alpha_1) \hat{g}_2(\alpha_1,\alpha_2) \ldots \hat{g}_d(\alpha_{d-1})$$

$$\hat{g}_k = \sum_{i_k} g_k$$

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Tensor Train and quadratures

$$I(d) = \int \sin(x_1 + x_2 + \ldots + x_d) \, dx_1 dx_2 \ldots dx_d =$$
$$\operatorname{Im} \int_{[0,1]^d} e^{i(x_1 + x_2 + \ldots + x_d)} \, dx_1 dx_2 \ldots dx_d = \operatorname{Im}((\frac{e^i - 1}{i})^d).$$

n = 11 nodes per one dimension \Rightarrow as amny as n^d values! But only neglible part is computed.

d	<i>I(d)</i>	Rel.error	Time
100	-3.926795e-03	2.915654e-13	0.77
1000	-2.637513e-19	3.482065e-11	11.60
2000	2.628834e-37	8.905594e-12	33.05
4000	9.400335e-74	2.284085e-10	105.49

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Tensorization of vectors and matrices

Any vector of size $N = n_1 \dots n_d$ can be regarded as a *d*-array, and any $N \times N$ matrix

$$a(i,j) = a(i_1 \dots i_d, j_1 \dots j_d)$$

can be viewed as as 2*d*-array, or better, as a *d*-array of size $n_1^2 \times \ldots \times n_d^2$. of the form

 $a(i_1j_1,\ldots,i_dj_d)$

Tensorization followed up by Tensor Train construction can dramatically decrease the number of representation parameters in the model!

We may have a blessing of dimensionality!

- For astronomically large data we must look for an adequate low-parametric models.
- For conventional big data tensorization may lead to tremendously efficient representation models.

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Thanks for your kind attention!

Эта книга отличается от традиционных учебников: понятия и факты линейной алгебры во многих случая изложены как факты матричного анализа Такий совтов релокат релокии соверсо	Д ш ш ш ш ш ш ш ш ш ш ш ш ш
абстрактым и позволет гозножение менее абстрактым и позволет гознакомить читателя с матричным анализом как относительно самостоятельной дисциплиной.	МАТРИЧНЫЙ НАНАЛИЗ
Предмет линейной алгебры понимается в расширенном смысле, часто мы оказываемся на территории смежных дисциплин — математического анализа, в выбологи изм исторов и слочения,	алгебра
ваниистипельных методов и, колечно, общей алгебры. Отдельные места в книге содержат материал, который вообще нельзя найти в какик-любо учебниках и даже монографиях. В частности, рассматриваются мигоомерные массивы,	илиней
классические и относительно новые тензорные разложения, тензорные ранги и их связь с разработкой быстрых методов умножения матриц.	
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Eugene Tyrtyshnikov

Mathematics and Computations